

METHOD OF BUFFER MANAGEMENT AND TASK SCHEDULING FOR
TWO-DIMENSIONAL DATA TRANSFORMING

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CROSS-REFERENCE TO RELATED APPLICATION

This application claims the priority benefit of provisional application serial no. 60/217,825, filed July 12, 2000, the full disclosure of which is incorporated herein by reference.

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BACKGROUND OF THE INVENTION

Field of Invention

The present invention relates to a method of buffer management and task scheduling. More particularly, the present invention relates to a method of buffer management and task scheduling for two-dimensional data transforming.

15 Description of Related Art

Two-dimensional Discrete Fourier Transforms, such as Discrete Cosine Transform (DCT), are widely used in image processing, digital signal processing, digital communications, etc. Due to its intrinsic properties, two-dimensional transform operation requires significant amount of memory to store data. One good example is a 20 two-dimensional DCT (2D-DCT) operated on a two-dimensional image. As illustrated in FIG. 1, an image 10 is scanned line-by-line from the left-top corner to the right-bottom corner. In order to perform 2D-DCT with 6x6 block size 12, at least five lines of data 16 have to be stored. Furthermore, new input data (not shown) keeps coming in while old data 14 are being processed. Since the input data follow a line-by-line pattern

and the old data 14 are processed in a block-by-block pattern, there is a conflict in memory space.

Traditionally, there are two solutions to solve this problem. The first solution is to use two six-line buffers, one for reading out data and the other for writing data in, and 5 then swapping the buffers every six lines. However, this requires double the memory space than required in theory. The second solution is to process the 2D-DCT very quickly and to clean out the memory space before the read-write conflict occurs. Though, this causes significant processing bursts while receiving the "last line" 18 of the 6x6 blocks 12. It does not only require very extensive computations for such 10 processing bursts, but also waste the computing resources while staying idle and waiting for all necessary data to be processed.

The invention provides a new solution to the problem stated above. The invention uses the minimum space required in theory and distributes the process loading evenly in time domain. It eliminates the trade-off between memory size and computing 15 power, which has been recognized as a serious dilemma in 2D Discrete Fourier Transforms. The invention also presents a practical implementation of these presented ideas which is very simple and, therefore, consumes much less power than traditional designs do.

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SUMMARY OF THE INVENTION

The invention provides a method of buffer management and task scheduling for two-dimensional data transforming, which transforms data in a buffer sequentially by blocks, wherein a block comprises a specific size having plurality of row portions and a plurality of column portions. The method comprises the steps of initializing a write

logic address and a read logic address to zero. Then, old data is read out in a block-by-block pattern. While reading out data, new data is written in immediately in a line-by-line pattern in the block using a first mapping scheme. After the whole block is read out, the block is moved to process another column portion. Then old data in the block is again read out in a block-by-block pattern while new data is again written in immediately in a line-by-line pattern in the block using a second mapping scheme. After the whole block is read out again, the block is moved to process another column portion. These read and write steps with the interleaving of the first and second mapping schemes are repeated until the whole data in the buffer is processed.

Prior to the read and write steps, the block is completely filled with data using a second mapping scheme. After the block is filled, the write logic address is incremented by 1. Accordingly, the write logic address is incremented by 1 per buffer write, and the read logic address is incremented by 1 per buffer read as well.

According to the invention, the first and second mapping scheme translates logical addresses to physical addresses, which are really used to access the buffer.

It is to be understood that both the foregoing general description and the following detailed description are exemplary, and are intended to provide further explanation of the invention as claimed.

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BRIEF DESCRIPTION OF THE DRAWINGS

The accompanying drawings are included to provide a further understanding of the invention, and are incorporated in and constitute a part of this specification. The drawings illustrate embodiments of the invention and, together with the description, serve to explain the principles of the invention. In the drawings,

FIG. 1 is a conventional method of buffer management and task scheduling for two-dimensional data transforming.

FIG. 2 is a logical addressing method of buffer management and task scheduling for two-dimensional data transforming in a preferred embodiment according to the
5 invention.

FIG. 3 is a flow chart overview of the operational procedure of buffer management and task scheduling for two-dimensional data transforming in a preferred embodiment according to the invention.

10 DESCRIPTION OF THE PREFERRED EMBODIMENTS

Reference will now be made in detail to the present preferred embodiments of the invention, examples of which are illustrated in the accompanying drawings. Wherever possible, the same reference numbers are used in the drawings and the description to refer to the same or like parts.

15 Before discussing the embodiments of the invention, a general and qualitative description of the invention should prove useful. The basic idea behind the invention is to keep the data input and output rates the same. In other words, in average, one datum from processing is taken out from the memory while writing one new datum into the same memory. By doing so, the memory space can be fully utilized and the process
20 loading can be evenly distributed. This may sound straightforward, however in reality, what just have been described is very hard to implement, since the input and output data follow different patterns. The 2D data is scanned in a line-by-line pattern, but they are processed in a block-by-block pattern. If new data is simply filled in the memory locations that are just emptied, the memory read-out pattern needs to be changed to

fulfill the original block-by-block output requirement. This means that the memory is being emptied in a different manner. Therefore, the pattern for writing the following data in the memory changes. In consequence, the memory read-out pattern for the next data changes again. The situation repeats itself. After a few iterations, the patterns for 5 writing and reading the memory look almost random. Unless a huge record to track all the previous input and output data is available, it is very hard, if not impossible, to know where to write the new data and read out the necessary data for processing. Keeping such a record is not practical because it will require much more memory space than needed in storing the original data. In this invention, by re-arranging the order of output 10 blocks, robust addressing schemes are created for input and output sequences to resolve the problem discussed above.

The operations of the novel buffer management and task scheduling for 2D-DCT is best described by an example. FIG. 2 is a logical addressing scheme of a buffer management and task scheduling for two-dimensional data transforming in a preferred 15 embodiment according to the invention.

In FIG. 2, assuming a 2D-DCT with 8x8 block size 22 is performed on an image 20 of 512 pixels per line. A 9-bit binary word 23 is used to address the pixels on each line, and a 3-bit binary word 25 is used to address the minimal 8 lines required for the 8x8 block size.

20 The 9-bit pixel address 23 is separated into 3-bit “dot”, 3-bit “block”, and 3-bit “sector” addresses. Associating this with the 3-bit line address 25, {3-bit line address, 9-bit pixel address} is equal to {3-bit line address, 3-bit sector address, 3-bit block address, 3-bit dot address}. Two logical addresses of this kind is obtained for buffer write and buffer read respectively. These two logical addresses are increased by one

from zero per access. In addition, two types of address mappings are defined for buffer write and buffer read. These two mappings translate logical addresses to physical addresses, which are really used to access the memory of the buffer. So, a write logical address, a write physical address, a read logical address, a read physical address, and
5 two address mapping schemes is obtained. The two address mapping schemes are listed as follows:

Mapping Scheme A: Physical Address = {Sector Address, Line Address, Block Address, Dot Address}

Mapping Scheme B: Physical Address = {Sector Address, Block Address, Line Address, Dot Address}

FIG. 3 is a flow chart overview of the operational procedure of buffer management and task scheduling for two-dimensional data transforming in a preferred embodiment according to the invention. In step 30, for the first time usage, the memory buffer is filled up using Mapping Scheme A for buffer writes. The write logic address
15 is increased by one from zero per buffer write.

Next, at step 31, when the buffer is full, it is ready to perform an 8x8-block size 2D-DCT operation. The read logic address is increased by one from zero per read write. Data of 8x8 block are read out by using Mapping Scheme B for buffer read. This mapping automatically guarantees the output sequence required by 2D-DCT operations.

20 Then, at step 32, while reading out data for 2D-DCT, the memory buffer is still receiving input data. The “one out, one in” method is followed to maximize the buffer efficiency. The input data by using Mapping Scheme B for buffer writes is then recorded.

Further, at step 33, after the whole buffer is read out for processing and is replaced by new data, it is ready to process another 8 lines stored in the buffer. It is essential to make sure the output sequence is correct for 2D-DCT operations. Data is read out by using Mapping Scheme A for buffer reads this time. This mapping
5 automatically guarantees the output sequence required by 2D-DCT operations.

Moreover, at step 34, as in step 32, the memory buffer is still receiving input data while reading out data for the 2D-DCT. The input data is recorded by using Mapping Scheme A for buffer writes.

Finally, after the whole buffer is read out for processing and is replaced by new
10 data, Mapping Scheme B is switched for both buffer reads and writes. As from step 31 to step 34, the two mapping schemes are interleaved to guarantee correct output sequences while the buffer is kept full all the time.

By interleaving the simple address mapping schemes, the dilemma in 2D data transforms can be resolved. In addition, the output flow is continuous after the first
15 filling. The task scheduling automatically follows the data output throughout. The process loading is smoothed out without loading bursts.

In the invention, a pre-buffer can be used to store data. This pre-buffer can be replaced by a “staging buffer” which stores temporary data between the “first one-dimensional” and “second one-dimensional” operations in 2D-DCT processes.. The
20 same concept can be applied to the staging buffer. In addition, the mapping schemes mentioned above can be modified or combined with more address manipulations to accommodate different operational requirements.

The invention achieves maximum buffer efficiency without trading off loading requirement. It does need a complicated address recording scheme to fulfill the output

sequences required by 2D-DCT operations. The 2D-DCT operations can also be scheduled accordingly. No existing methods can resolve the same dilemma without increasing the system complexity significantly.

It will be apparent to those skilled in the art that various modifications and
5 variations can be made to the structure of the present invention without departing from the scope or spirit of the invention. In view of the foregoing, it is intended that the present invention cover modifications and variations of this invention provided they fall within the scope of the following claims and their equivalents.